

Learning and Robustifying an Optimal Assortment from Observational Data

Miao Lu

Stanford University



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Joint work with



Yuxuan Han
NYU



Han Zhong
PKU



Zhengyuan Zhou
NYU



Jose Blanchet
Stanford

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Assortment Optimization: Math. Setups

Seller (assortment and revenue)

Seller wants to maximize expected revenue.

- ▶ Product set: $[N] := \{1, 2, \dots, N\}$.
- ▶ An assortment: $S \subseteq [N]$.
- ▶ Revenue of product i : $r_i \in [0, r_{\max}]$.
- ▶ Non-purchasing: $i = 0$ and $r_0 = 0$.

Customer (choice model)

Customer buy product based on choice model:

- ▶ $\mathbb{P}(i|S) = \mathbb{P}(\text{buying } i \text{ seeing assortment } S)$.
- ▶ This work: multinomial logit (MNL),
$$\mathbb{P}_{\mathbf{v}}(i|S) := \frac{v_i}{1 + \sum_{j \in S} v_j}, \quad \forall i \in S \cup \{0\}.$$

Expected Revenue of assortment S :

$$R(S; \mathbf{v}, \mathbf{r}) := \sum_{i \in S} r_i \cdot \mathbb{P}_{\mathbf{v}}(i|S) = \sum_{i \in S} \frac{r_i v_i}{1 + \sum_{j \in S} v_j}.$$

Optimal Assortment S^* that maximizes the expected revenue:

$$S^* := \arg \max_{|S| \leq K} R(S; \mathbf{v}, \mathbf{r}), \quad K \in \mathbb{N}: \text{ size constraint.}$$

Learning Optimal Assortment S^* from Data

This Work: Offline Learning of S^* .

- ▶ The seller doesn't know the underlying choice pattern $\mathbb{P}_{\mathbf{v}}$, but still wants to present good S .
- ▶ Learning S^* purely from offline data of customer choices, denoted by $\mathbb{D} := \{(S_k, i_k)\}_{k=1}^n$.
- ▶ $S_k \subseteq [N]$ is an observed assortment, $i_k \sim \mathbb{P}_{\mathbf{v}}(\cdot | S_k)$ is the observed customer choice.
- ▶ Performance metric (sub-optimality gap): $\text{SubOpt}(\hat{S}; \mathbf{v}, \mathbf{r}) := R(S^*; \mathbf{v}, \mathbf{r}) - R(\hat{S}; \mathbf{v}, \mathbf{r})$.

Key Questions and Result Overview:

1. (Lower bound) What is the minimal condition on \mathbb{D} for learning S^* under the MNL model?
A new notion of "item-wise partial coverage" condition.
2. (Upper bound) What is the algorithm to learn S^* requiring only the minimal condition on \mathbb{D} ?
Pessimistic Rank-Breaking (PRB), achieving minimax optimal sample complexity.
3. (Robustness) What if the observational data mismatches the actual customer choice pattern?
Extensions of 1. & 2. to robust assortment optimization.

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What is the Minimal Condition on Data \mathbb{D} ?

- ▶ Arbitrarily bad offline data can not guarantee the seller to learn a good assortment!
- ▶ “Partial coverage” is the de facto condition for effective offline data-driven decision making: the data needs and only needs to cover the optimal action.
 - In contrast with “full coverage” style conditions, which require uniform coverage of action space.
- ▶ In assortment optimization, the optimal action can be viewed as S^* .
- ▶ Prior work [Dong et al. \[2023\]](#) proposes PASTA, which achieves efficient learning if S^* appears sufficiently often in \mathbb{D} : $\text{SubOpt}(\hat{S}) \lesssim 1/\sqrt{n_{S^*}}$, where $n_{S^*} := \sum_{k=1}^{|\mathbb{D}|} \mathbf{1}\{S_k = S^*\}$.
- ▶ However, the vanilla “partial coverage” condition could still be impractical here!
- ▶ Combinatorial nature of assortment S : if we randomly sample items to form an assortment of size K , the prob. of observing S^* once would be of order $\Theta(N^{-K})$.
- ▶ Can we further relax the vanilla partial coverage condition?
 - This work: “item-wise partial coverage condition” is both necessary and sufficient!

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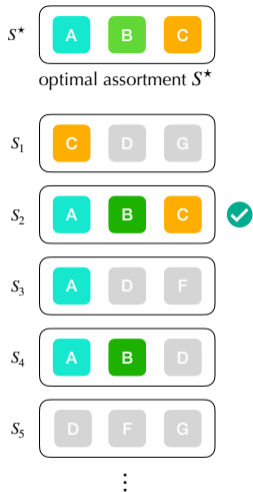
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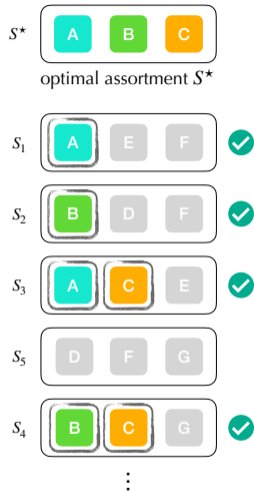
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Data Pattern needed for previous SOTA
(Require sufficient S^* in data)



Our algorithm learns S^* from this data pattern
(No S^* appearing in the data)

Our Theoretical Results: An Overview

Item-wise Partial Coverage (Ours). Effective size of \mathbb{D} :

$$n_{\text{effect}} := \min_{i \in S^*} n_i, \quad n_i := \sum_{k=1}^{|\mathbb{D}|} \mathbf{1}\{i \in S_k\}.$$

Vanilla Partial Coverage [Dong et al., 2023].

$$n_{S^*} := \sum_{k=1}^{|\mathbb{D}|} \mathbf{1}\{S_k = S^*\} \leq n_{\text{effect}}.$$

	Setup	Upper Bound	Lower Bound
[Dong et al., 2023]	general revenue	$\tilde{O}\left(\sqrt{\frac{NK}{n_{S^*}}}\right)$	N/A
This Work	general revenue	$\tilde{O}\left(\frac{K}{\sqrt{n_{\text{effect}}}}\right)$	$\Omega\left(\frac{K}{\sqrt{n_{\text{effect}}}}\right)$
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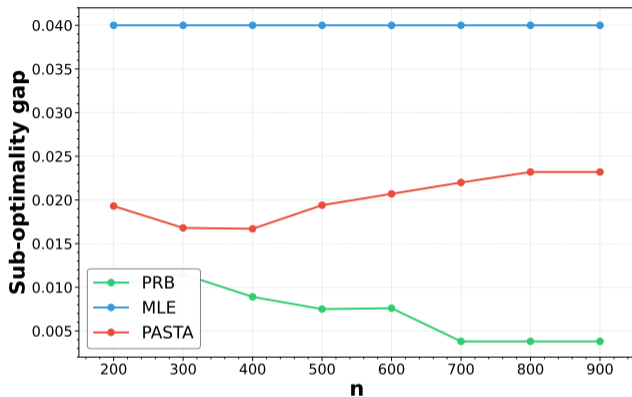
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- ▶ PRB: Our algorithm
- ▶ PASTA [Dong et al., 2023], MLE (naive baseline).

Our Algorithm: Pessimistic Rank-Breaking (PRB)

Takeaways

- ▶ For offline A.O. with MNL model, “item-wise partial coverage” is the minimal data condition.
- ▶ PRB achieves minimax-optimal sample complexity, requiring only item-wise coverage.
- ▶ There is a \sqrt{K} -statistical gap in general revenue case and the uniform revenue case.

Robust Assortment Optimization

- ▶ So far we assumed that the choice pattern \mathbb{P} generating \mathbb{D} matches the future customer choice distribution.
- ▶ What if the learning from \mathbb{P} generates a wrong choice model for the future?
- ▶ Our solution: robust assortment optimization from observational data.

Definition (Optimal Robust Assortment)

Given nominal choice model \mathbb{P} , robust level function ρ , we define optimal robust assortment as

$$S_{\rho(\cdot; \cdot)}^* = \arg \max_{|S| \leq K} R_{\rho(\cdot; \cdot)}(S; \mathbb{P}, \mathbf{r}), \quad (1)$$

$$R_{\rho(\cdot; \cdot)}(S; \mathbb{P}, \mathbf{r}) := \inf_{\substack{Q_{S_+} \in \mathcal{P}(S_+) \\ D_{\text{KL}}(Q_{S_+} \| \mathbb{P}(\cdot | S)) \leq \rho(S; \mathbb{P}(\cdot | S))}} \sum_{j \in S_+} Q_{S_+}(j) \cdot r_j. \quad (2)$$

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Robust Item-wise Partial Coverage. Effective size of \mathbb{D} :

$$n_{\text{effect};\rho(\cdot;\cdot)} := \min_{i \in S_{\rho(\cdot;\cdot)}^*} n_i, \quad n_i := \sum_{k=1}^{|\mathbb{D}|} \mathbf{1}\{i \in S_k\}.$$

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Takeaways

- ▶ For offline robust A.O. with MNL as nominal choice model, “robust item-wise partial coverage” is the minimal data condition.
- ▶ New algorithms to achieve minimax-optimal sample complexity.
- ▶ There keeps to be a \sqrt{K} -statistical gap in general revenue case and the uniform revenue case.

Thank you for your attention!

References I

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Lower Bound

Theorem (Minimax Lower Bound 1 [Han et al., 2025])

For sufficiently large N, K with $N \geq 5K$, there is an \mathbb{D} of K -sized assort. $\{S_k\}_{k=1}^n$ and a class of problem instances \mathcal{V} s.t.,

$$\inf_{\pi(\cdot): \mathbb{D} \mapsto S} \sup_{(\mathbf{v}, \mathbf{r}) \in \mathcal{V}} \mathbb{E}_{\mathbb{D} \sim \otimes_{k=1}^n \mathbb{P}_{\mathbf{v}}(\cdot | S_k)} [\text{SupOpt}(\pi(\mathbb{D}); \mathbf{v}, \mathbf{r})] \geq \Omega \left(\frac{r_{\max} \cdot K}{\sqrt{n_{\text{effect}}}} \right). \quad (3)$$

Corollary (noframenumbering)

[Minimal Data Condition: Item-wise Partial Coverage] The minimal condition on the observational data \mathbb{D} is to have a good coverage of each single product i in the optimal assortment S^* .

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Our Algorithm: Pessimistic Rank-Breaking (PRB)

Two key steps: (i) rank-breaking based estimation; (ii) pessimistic assortment optimization.

Step 1: Rank-breaking Estimation [Saha and Gopalan, 2019, Saha and Gaillard, 2024].

$$\hat{p}_j = \frac{\tau_j}{\tau_{j0}}, \quad \text{where} \quad \tau_{j0} = \underbrace{\sum_{k=1}^{|\mathbb{D}|} \mathbf{1}\{i_k \in \{0, j\}, j \in S_k\}}_{\# \text{ Choosing item } j \text{ or leave}}, \quad \tau_j = \underbrace{\sum_{k=1}^{|\mathbb{D}|} \mathbf{1}\{i_k = j\}}_{\# \text{ Choosing item } j}, \quad (4)$$

- ▶ Insight: pairwise choice probabilities $p_j := \mathbb{P}(j \mid S = \{0, j\}) = v_j / (1 + v_j)$ for different $j \in [N]$ can be estimated separately, which can then be mapped to v_j , i.e., $\hat{v}_j = \hat{p}_j / (1 - \hat{p}_j)$.

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Step 2: Pessimistic Assortment Optimization.

To hedge the potential large uncertainty for those items not well covered by \mathbb{D} , we optimize the assortment against a conservatively estimated revenue. Firstly,

$$p_j^{\text{LCB}} = \left(\hat{p}_j - \sqrt{\frac{2\hat{p}_j(1-\hat{p}_j)\log(1/\delta)}{\tau_{j0}}} - \frac{\log(1/\delta)}{\tau_{j0}} \right)_+ \leq p_j \quad (5)$$

Then we optimize the assortment to maximize the pessimistic estimation of the expected revenue:

$$\hat{S} := \arg \max_{|S| \leq K} R(S; \mathbf{v}^{\text{LCB}}, \mathbf{r}), \quad v_j^{\text{LCB}} = p_j^{\text{LCB}} / (1 - p_j^{\text{LCB}}) \leq v_j, \quad \forall j \in [N], \quad (6)$$

Lemma (Monotonicity under Optimal Assortment [Han et al., 2025])

Given $\{\underline{v}_j\}_{j=1}^N$ s.t. $v_j \geq \underline{v}_j, \forall j \in [N]$. Then for the maximizer $\underline{S} := \arg \max_{|S| \leq K} R(S; \underline{\mathbf{v}}, \mathbf{r})$,

$$R(S^*; \mathbf{v}, \mathbf{r}) - R(\underline{S}; \mathbf{v}, \mathbf{r}) \leq \sum_{j \in S^*} \frac{r_j(v_j - \underline{v}_j)}{1 + \sum_{j \in S^*} v_j}. \quad (7)$$

Our Algorithm: Pessimistic Rank-Breaking (PRB)

Theorem (Sub-Optimality Gap of PRB [Han et al., 2025])

Assume that $n_i \geq 256(1 + KV) \log(N/\delta)$ for all $i \in S^*$. Then, with probability at least $1 - 2\delta$, the output \hat{S} of PRB satisfies that

$$\text{SubOpt}(\hat{S}; \mathbf{v}, \mathbf{r}) \leq 2 \cdot r_{\max} \cdot (1 + v_{\max}) \cdot K \cdot \sqrt{\frac{\log(N/\delta)}{n_{\text{effect}}}} + \text{higher order terms.} \quad (8)$$